

## Math 4021 – Prof. Pennance – On Substitution

1. Let  $X$  a set variables,  $F$  a formula and  $q : X \rightarrow \text{Term}$  a “term assignment”. Our aim is twofold.

- (a) To define  $F(q)$  – the formula obtained by “substituting” each  $x \in V_f(F) \cap X$  by  $q(x)$ .
- (b) To define carefully when this substitution makes sense.

2. Definition.

- (a) If  $c$  is a constant  $c(q) = c$
- (b) If  $x$  is a variable

$$x(q) = \begin{cases} q(x), & \text{if } x \in X; \\ x, & \text{if } x \notin X; \end{cases}$$

3.  $F(q)$  is now defined recursively as follows:

- (a)  $(t_1 = t_2)(q) = t_1(q) = t_2(q)$
- (b)  $rt_1 \dots t_n(q) = rt_1(q) \dots t_n(q)$  where  $r$  is a relational symbol of order  $n$ .
- (c)  $(\neg F)(q) = \neg F(q)$
- (d)  $(F \rightarrow G)(q) = F(q) \rightarrow G(q)$
- (e)  $(\forall x F)(q) = \forall x F(q')$  where  $q'$  is the restriction of  $q$  to  $X' = X - x$

4. If  $q(x_1) = t_1, q(x_2) = t_2, \dots, q(x_n) = t_n$  then  $F(q)$  is often denoted  $F \left[ \begin{smallmatrix} x_1 \dots x_n \\ t_1 \dots t_n \end{smallmatrix} \right]$

In the special case  $n = 1$ , in which a single variable  $x$  is substituted by a term  $t$ , definition (3) reads:

- (a) If  $F$  atomic then  $F \left[ \begin{smallmatrix} x \\ t \end{smallmatrix} \right]$  is obtained from  $F$  by replacing  $x$  by  $t$ .
- (b)  $(\neg F) \left[ \begin{smallmatrix} x \\ t \end{smallmatrix} \right] = \neg F \left[ \begin{smallmatrix} x \\ t \end{smallmatrix} \right]$

$$(c) (F \rightarrow G) \left[ \begin{smallmatrix} x \\ t \end{smallmatrix} \right] = F \left[ \begin{smallmatrix} x \\ t \end{smallmatrix} \right] \rightarrow G \left[ \begin{smallmatrix} x \\ t \end{smallmatrix} \right]$$

$$(d) (\forall x F) \left[ \begin{smallmatrix} y \\ t \end{smallmatrix} \right] = \begin{cases} \forall x F & \text{if } y = x; \\ \forall x F \left[ \begin{smallmatrix} y \\ t \end{smallmatrix} \right] & \text{if } y \neq x; \end{cases}$$

5. It follows from (3) that only free variables are subject to change under a substitution. For example,

$$(\forall x x+y = y+x) \left[ \begin{smallmatrix} x \\ z \end{smallmatrix} \right] = \forall x x+y = y+x$$

6. Warning. If we fail to restrict  $q$  and use instead the INCORRECT formula

$$(\forall x F)(q) = \forall x F(q)$$

we obtain the INCORRECT result

$$(\forall x x+y = y+x) \left[ \begin{smallmatrix} x \\ z \end{smallmatrix} \right] = \forall x z+y = y+z$$

which does not have the intended meaning.

7. There is still a problem to be overcome. For example, using the correct substitution rule (3e) we have

$$(\forall x x+y = y+x) \left[ \begin{smallmatrix} y \\ x \end{smallmatrix} \right] = \forall x x+x = x+x$$

which is incorrect semantically. The problem arises because, even though  $y$  is free, it falls within the scope of the quantifier  $\forall x$ . It follows that any term  $t$  containing the variable  $x$ , (in the above example  $t = x$ ) will be subject to quantification upon substitution for  $y$ . We must therefore limit ourselves to substitutions which are “proper” for logical deduction. The notation of “proper” is made precise by the following definition.

8. Let  $q$  be a term assignment and  $F$  a formula. We define the relation  $S(q, F)$  read  $q$  is a *proper substitution* in  $F$  as follows:

- (a) If  $F$  is atomic, then  $S(q, F)$ .
- (b) If  $S(q, F)$  then  $S(q, \neg F)$ .
- (c) If  $S(q, F)$  and  $S(q, G)$  then  $S(q, F \rightarrow G)$
- (d)  $S(q, \forall x F)$  if
  - i.  $S(q', F)$  and
  - ii. For all  $v \in V_f[\forall x F \cap X]$ ,  $x \notin V(q(v))$ .  
(recall,  $V(t)$  denotes the set of variables of a term  $t$ )

Note: (ii) can be written

$$x \notin V(q[V_f(\forall x F)])$$

Thus a substitution  $q$  in  $\forall x F$  is improper if either  $q'$  is improper in  $F$ , or, for some free variable  $v \in \forall x F$ ,  $q(v) = t$  where  $t$  contains the quantified variable  $x$ .

## 9. Examples

- (a)
 
$$(\exists x x + x = y) \left[ \frac{y}{x} \right] = \exists x x + x = x$$

is improper since  $V(q(y)) = \{x\}$  and  $x$  is quantified.

(b)

$$(\forall x x + y = y + x) \left[ \frac{x}{y} \right]$$

is proper. In this case, the domain of  $q'$  is empty and no free variables is eligible for substitution by a term. Thus, the substitution yields

$$\forall x x + y = y + x$$

(c)

$$(\forall y \forall x x + y = y + x) \left[ \frac{x}{y} \right]$$

equals

$$\forall y (\forall x x + y = y + x) \left[ \frac{x}{y} \right]$$

which by the previous example equals

$$\forall y \forall x x + y = y + x$$

Notice that the substitution leaves the formula unchanged. It is left as an exercise to show that the substitution is proper.